Card-based Cryptographic Protocols for Three-input Functions with a Standard Deck of Cards Using Private Operations

Naoki Kobayashi <u>Yoshifumi Manabe</u>
Kogakuin University
Shinjku, Tokyo Japan

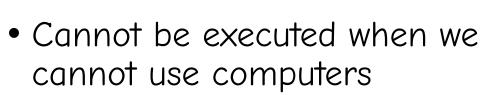


Outline of Talk

- What is "Card-based Cryptographic protocols?"
- Private operations
- Problem in using standard deck of cards
- Proposed protocols
- Conclusion

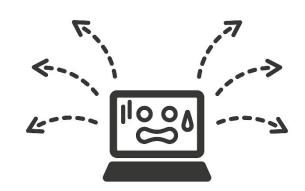
Cryptographic protocols using computers

 Difficult to understand the correctness for non-experts



 Software might not be reliable: Is the private data really discarded?



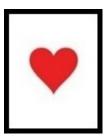


Card-based Protocols

Cryptographic

• Using cards





Back of the cards



- Can be executed when we cannot use computers
- Relatively easy to understand the protocol →can be used for teaching cryptography
- Private data are lost after the cards are shuffled



Overview of card-based cryptographic protocols

- Two kinds of protocols
 - Primitives to calculate any functions
 - AND, XOR, copy of inputs
 - Efficient protocols to solve specific problems
 - Voting/Grouping/Millionaires' Problem/etc.
- Executed by semi-honest Alice and Bob
 - Obey the rule but try to obtain private values
 - Private values must not be known to the players
- Most protocols use special cards
 - ⇒Use standard deck of cards

Special cards used by most protocols





Back is



indistinguishable

Encoding of one bit data





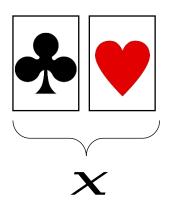
= 1



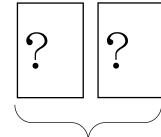


= C

• Data and its commitment



Face down



commit(x)

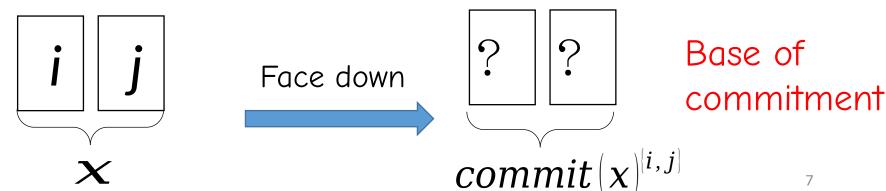
(This paper) Standard deck of cards

1 | 2 | ... | 52 | Back is | ? indistinguishable

Encoding of one bit data

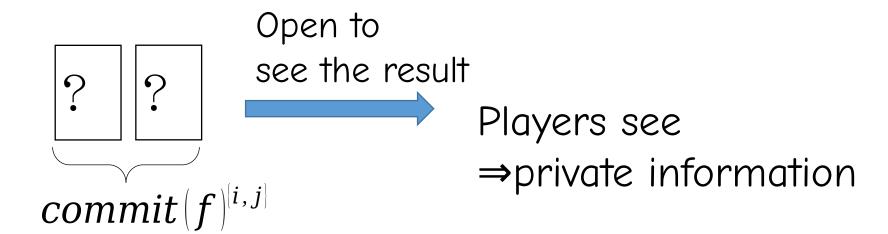
$$|j|$$
 $|i|$ =1 $|i|$ $|j|$ =0 if $i < j$

Data and its commitment



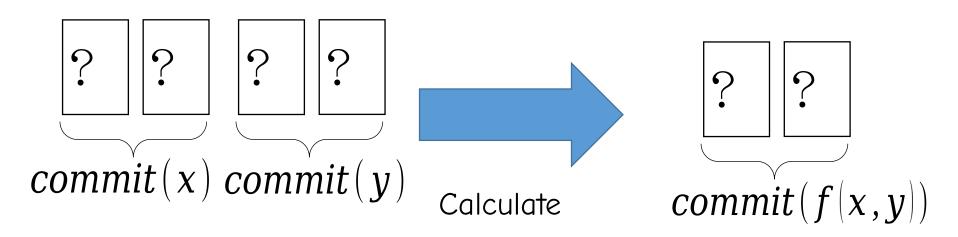
Difference between — cards and standard deck of cards

Information leakage from the bases of the cards



Requirements for inputs and outputs

- Input: protocols must arrow committed inputs
 - Players can calculate using unknown values
- Output: protocols must output committed results
 - Can be used as inputs to further computation



private operations

- Some player executes operations where the other players cannot see (under the table, in the back): private operations
- [NTMI016] Millionaires' problem
- [OM18] : Minimum number of cards for AND, XOR, Copy protocols

Number of additional cards (uniform closed shuffle, finite step)

	cards	Without Private operation	With Private operation
AND		2 [MS09]	O [OM18]
	standard	4 [M16]	O [MO24]
XOR	•••	0 [MS09]	0 [OM18]
	standard	O [M16]	O [MO24]
Сору	•••	2 [MS09]	0 [OM18]
	standard	2 [M16]	0 [MO241]

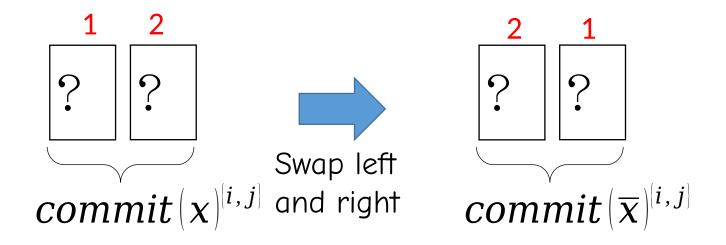
Number of additional cards (uniform closed shuffle, finite step)

	cards	Without Private operation	With Private operation
Three input fn.		2 [NHMS15]	0 [MO21]
	standard	4 [M16]+↑	O [This]
Half/Full adder		2 [NHMS15]	0 [MO21]
	standard	4 [M16]+1	O [This]
Symmetri c fn.	•••	2 [NHMS15]	0 [MO21]
	standard	4 [M16]+↑	O [This]

Primitives used in the protocols

- NOT
- AND with private operations
- (XOR, copy with private operations)
- XOR, AND with preserving an input
- Any Boolean function with 4 additional cards

Primitive executable by anyone



AND protocol[OM18]

• Input: , Output:

(1)Alice(private): Random swap Randomly select $commit(x \oplus a_1)$ commit(x)(2)Bob(private): OR See →Set cards commit(y) commit(0) commit(0) commit(v) $x \oplus a_1 = 1 \times a_1 = 0$

(3)Alice(private):

select left pair if right pair if

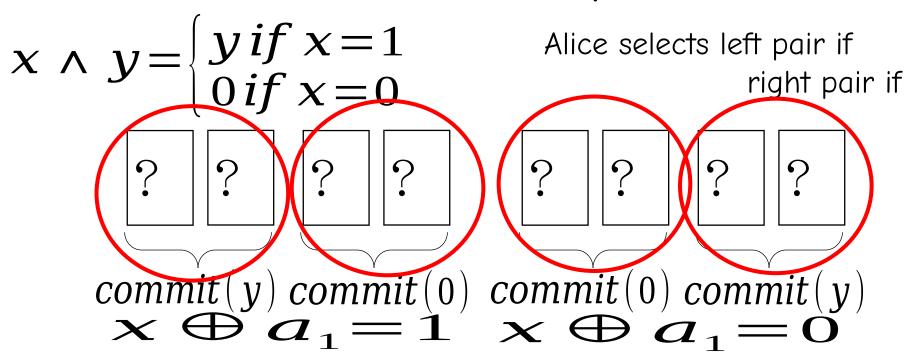
of cards: 4

re-use

to set

15

Correctness of AND protocol



Selects : and) or

Selects: and) or

(and)
$$\longrightarrow x=1$$

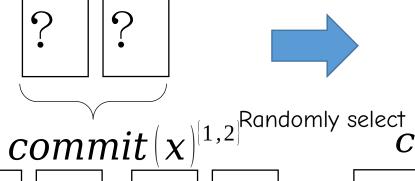
$$(and) \qquad \longleftarrow \qquad x = 0$$

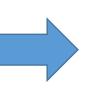
AND:standard cards $x \wedge y = \begin{vmatrix} y & if & x = 0 \\ 0 & if & x = 0 \end{vmatrix}$

• Input:

(1)Alice: Random swap







OR





 $Commit(x \oplus a_1)^{|1|}$

(2)Bob:









commit



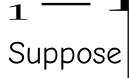




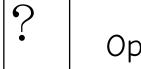
(3)Alice:

Select left pair if right pair if

 $commit(y)^{[3,4]}commit(0)^{[1]}commit(0)^{[1,2]}commit(y)^{[3,4]}$ = 1







Open cards

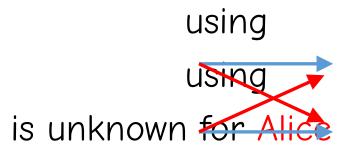
→sees

17

For execution using standard cards

Prevent private input data leakage from the bases of cards

• (Definition) semi-opaque commitment pair



AND protocol[MO24]

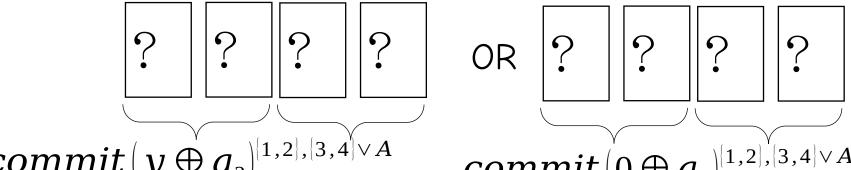
• Input: , Output:

(1)Alice(private): Random swap $commit(x)^{[1,2]}$ Randomly select $commit(x \oplus a_1)^{[1,2]}$ (2)Bob(private) See →Set cards $commit(y)^{[3,2]}commit(0)^{[1,2]}commit(0)^{[1]}commit(y)^{[3,2]}commit$

AND protocol(2)

(3) Alice: random ?????OR???????? Swap $commit(y \oplus a_2)^{[3,4]} commit(0 \oplus a_3)^{[1,2]} commit(y \oplus a_3)^{[3,4]} commit(y \oplus$

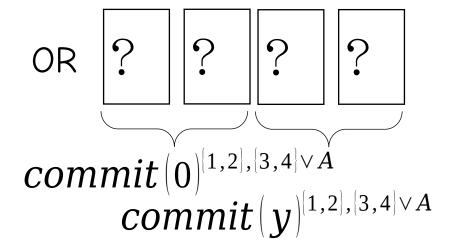
Randomly select and swap $\{1,2\}\{3,4\}$ if



 $commit(y \oplus a_2)^{[1,2],[3,4]\vee A} commit(0 \oplus a_2)^{[1,2],[3,4]\vee A} commit(y \oplus a_3)^{[1,2],[3,4]\vee A} commit(y \oplus a_3)^{[1,2],[3,4]\vee A}$

AND protocol(3)

(5) Alice: ? ? ? ? ? ? swap $commit(y)^{[1,2],[3,4]\vee A}$ $commit(0)^{[1,2],[3,4]\vee A}$

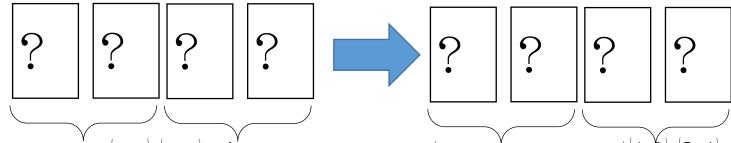


unselected pair:

AND protocol(4)

(6) Bob: For the selected pair and unused pair

Random swap using



 $commit(x \land y)^{[1,2],[3,4]\lor A^{'}}$ unselected pair:

$$commit(x \land y \oplus br_1)^{[1,2],[3,4] \lor A}$$

$$commit(\bot \oplus br_2)^{[1,2],[3,4] \lor A}$$

(7) Alice: see cards and set using $\{1,2\}$



(8) Bob: undo randomization by →obtain

 $commit(\overset{\mathsf{Y}}{x} \wedge y \oplus br_1)^{[1,2]} commit(\overset{\mathsf{Y}}{x} \wedge y)^{[1,2]}$

XOR with preserving an input

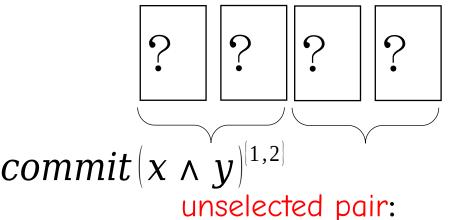
• Input : Output : $commit(x)^{[1,1]}$ (1) Alice: Random swap $commit(x)^{[1,2]}$ Randomly select $\widetilde{commit}(x \oplus b)^{[1,2]}$ (2) Bob: (Private) OR see: →set cards returns $commit(\overline{\mathbf{y}})^{[3,4]}$ $commit(y)^{[3,4]}$ $x \oplus b = 1$ $x \oplus b = 0$ \leftarrow $commit(y \oplus (x \oplus b))^{[3,4]}$ (3) Alice: Undo Swap if $\overline{\operatorname{commit}}(y \oplus (x \oplus b) \oplus b)^{[3,4]} = \operatorname{commit}(x \oplus y)^{[3,4]}$ $commit(x)^{[1,2]}$ 23

AND with preserving an input

• Input : Output :

 $commit(y)^{[3,4]}$

For the selected pair and unused pair



Unselected pair:

XOR with preserving one input between and

 \Rightarrow

Any Boolean Function with 4 additional cards[MO24]

```
    Two pairs(initial value 0)

    : Store intermediate value

    Boolean function can be written

For to
   copy
   For to
         (AND preserving input: )
```

Outline of the result

- Three-input Boolean function
- Half and full adder
- Any symmetric Boolean function

From the next page, is written as

Three-input Boolean functions

Categorized to the following 14 functions by considering swapping input variables and negations[SB10]



$$NPN_7 = |x \wedge y| \vee |x \wedge z| = x \wedge |y \vee z|$$

$$NPN_{13} = |\chi \wedge y \wedge z| \vee |\chi \wedge \overline{y} \wedge \overline{z}| = \chi \wedge |\overline{y} \oplus z|$$

$$NPN_6 = |x \wedge y \wedge z| \vee |\overline{x} \wedge \overline{y} \wedge \overline{z}|$$

- Input
 - XOR preserving NOT
 - XOR,NOT

AND

:variables current operation is applied

Red: newly obtained value

$$NPN_8 = |x \wedge y| \vee |\overline{x} \wedge \overline{y} \wedge z|$$

- Input
 - XOR with preserving ,NOT
 - OR(using AND)
- ,
 - **AND**

$NPN_9 = (x \wedge y \wedge \overline{z}) \vee$

- Input

XOR with preserving input

XOR with preserving input NOT



OR



$$NPN_{12} = |x \wedge \overline{z}| \vee |y \wedge z|$$

• Input

AND with preserving input

AND

OR

$$NPN_{11} = |x \wedge y| \vee |x \wedge z| \vee (y \wedge z)$$

• Input



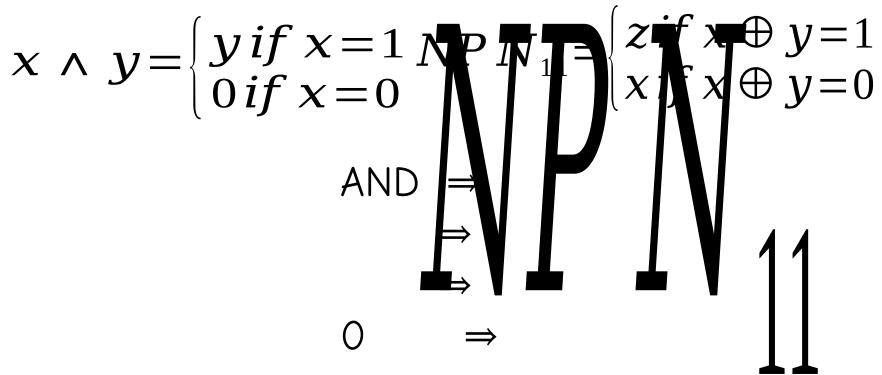
XOR with preserving input

Apply modified AND protocol



Modified AND protocol

AND



(2)Half adder Output:

• Input XOR with preserving input NOT AND with preserving input

Full adder Output: ,

• Input



Half adder



Half adder



OR(using AND and NOT)

(3) Symmetric function for any

is decided by the numbers of

$$k = l \log n J + 1$$

be binary representation



Half adder, Full adder



Realization of any Boolean function

(1)

Result is 3 bits \Rightarrow Any three-input can be realized without additional cards \Rightarrow can be realized without additional cards

(2)

Result bits

⇒unused input cards

can be realized with 4 additional cards

⇒ can be realized without additional cards

Conclusion

- Using private operations, minimum card protocol using standard deck of cards
 - Three-input Boolean function
 - Half and full adder
 - Symmetric Boolean function
- Open problems
 - -variable Boolean function
 - Minimizing the number of steps

Backup slides

$$x \land y = \begin{cases} y & \text{if } x = 0 \\ 0 & \text{if } x = 0 \end{cases}$$

Security of AND protocol

- Bob: sees →no information about because of randomization by
- Alice: sees ,

$$b_3 \equiv 0$$
 $b_2 \equiv 1$

Values are randomized by

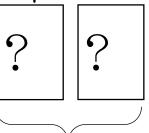
Sees {1,2} or {3,4} is randomized by

→No information from the cards

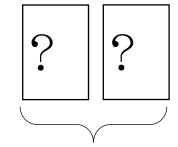
COPY protocol

- Input : , Output :
- (1) Alice:

Random swap





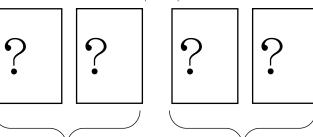


 $commit(x)^{[1,2]}$

 $commit(x \oplus b)^{[1,2]}$

- (2) Bob:
- See

Copy the value



 $commit(x \oplus b)^{[1,2]}commit(x \oplus b)^{[3,4]}$

- (3) Alice : Undo randomization for each pair

 $commit(x)^{[1,2]}commit(x)^{[3,4]}$

XOR protocol

• Input : Output :

(1) Alice: andom Randomly select wap $commit(x)^{[1]}$ $commit(y)^{[3,4]}$ $commit(x \oplus b)^{[1,2]}$ $commit(y \oplus b)^{[3,4]}$ (2) Bob: OR see and swap $commit(\overline{x \oplus b})^{[1,2]}commit(x \oplus b)^{[1,2]}$ if the value is 1 $v \oplus b = 1$ $v \oplus b = 0$

Correctness:
Output is

Any n-variable Boolean function

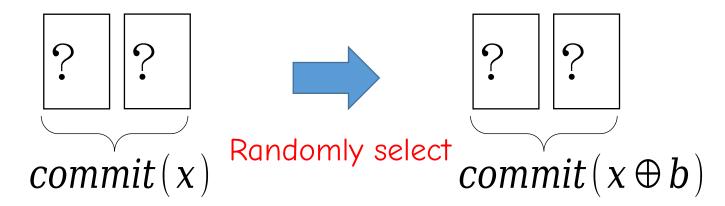
Input:

Output

- (1) 4 rounds, # of cards
- (2) rounds, # of cards

Private operation(1) [42]

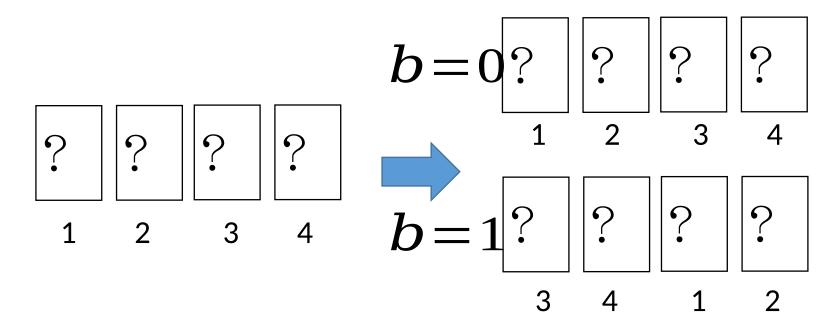
 Private random bisection cut: execute random bisection cut[31] in a hidden place



Remember. Do not disclose to any other player

Private operation(2) [42]

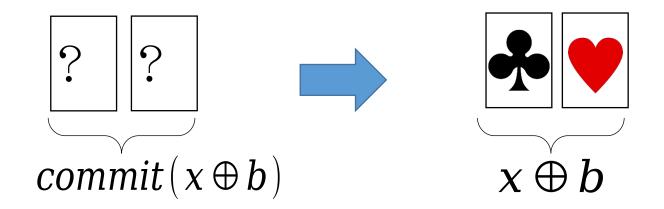
• Private reverse cut: swaps left and right of an even sequence using remembered



 (Private reverse selection : select left if right if)

Private operation(3) [42]

Private reveal: open cards and read the value



Executed by a player who does not know
→No information about

AND protocols

	Cards	# of cards	Note
[MS09]	*	6	
[RI18]	*	4	Las Vegas algorithm <i>†</i>
[OM18]	*	4	Use private operations
[NR99]	Standard	5	Las Vegas algorithm <i>†</i>
[KSK21]	Standard	4	Las Vegas algorithm <i>†</i>
[M16]	Standard	8	
[MO24]	Standard	4	Use private operations

[†] Bounded average execution time, no termination if random values are bad

Four cards are necessary for two inputs

Copy protocols

	Cards	# of cards	Note
[MS09]	♥♣	6	
[OM18]	♥♣	4	Use private operations
[NR99]	Standard	6	Las Vegas algorithm †
[M16]	Standard	6	
[MO24]	Standard	4	Use private operations

4 cards are necessary for copy

[†] Bounded average execution time, no termination if random values are bad

XOR protocols

	Cards	# of cards	Note
[MS09]	₩.	4	
[OM18]	••	4	Use private operations
[NR99]	Standard	4	Las Vegas algorithm +
[M16]	Standard	4	
[MO24]	Standard	4	Use private operations

† Bounded average execution time, no termination if random values are bad

Four cards are necessary for two inputs